

CSci 148, Spring 2002, Exam 2

Name (please print): \_\_\_\_\_ ID #: \_\_\_\_\_  
First Last

- This is a 75-minute closed-book exam. There are six questions worth a total of 95 points. Everybody gets 5 free points, so 100 points are possible. Budget your time so you get to all the questions. If a question seems unclear, write a short note about what assumption(s) you are making, and proceed as best as you can.

- |          |              |
|----------|--------------|
| 1. _____ | 5. _____     |
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| 4. _____ | Total. _____ |

(for grader use only)

**Question 1. (12 points)** Give a value of type `int event` that, when synchronized, does the following: (a) reads an integer value from channel `ch1`, then (b) reads an integer value from channel `ch2`, then (c) writes the sum of the two integers to channel `ch3`, and then finally (d) produces the (integer) sum. This event should be constructed in such a way that it has communication (b) as its commit point.

**Question 2. (10 points)** The book defines a function with signature

```
val mkEvent : 'a event * string * string -> 'a event
```

such that synchronizing on the event `mkEvent(e, ackMsg, nackMsg)` first synchronizes on `e`, and then if `e` is chosen, prints `ackMsg`, whereas if `e` is not chosen, prints `nackMsg`. For example, the synchronization

```
select [
  mkEvent(recvEvt ch1, "ch1\n", "not ch1\n"),
  mkEvent(recvEvt ch2, "ch2\n", "not ch2\n")
]
```

will either print the messages "ch1" and "not ch2" (in some order) or print the messages "ch2" and "not ch1" (in some order), depending on whether the communication on `ch1` or `ch2` is enabled first. Give a definition of `mkEvent`. (Hints: Use `withNack`, spawning a separate process in the "guard" function to wait for the negative acknowledgement. Printing can be achieved using `TextIO.print` of type `string -> unit`.)

**Question 3.** (23 = 16 + 6 + 3 points) Suppose that there are two global channels defined:

```
val reqCh : (int * int ivar) chan    (* request channel *)
val updCh : int chan                (* update channel *)
```

For this question, you will define three functions with signatures

```
val addServer : int -> unit    (* RPC server constructor *)
val addn      : int -> int     (* client call #1      *)
val update    : int -> unit    (* client call #2      *)
```

and with the following behavior:

- The call `addServer n` creates an “add  $n$ ” server, with local state  $n$ , that accepts requests on the two global channels. An *add* request of the form  $(i, iv)$  on `reqCh` stores the value of  $i+n$  into the I-variable  $iv$  and leaves the server state unchanged. An *update* request of the form  $m$  on `updCh` turns the server into an “add  $m$ ” server (i.e., its local state becomes  $m$ ).
- The call `addn i` computes and returns  $i + n$  by making an add request on `reqCh` to the server ( $n$  is the local state of the server)
- The call `update m` makes an update request on `updCh` to the server, turning it into an “add  $m$ ” server.

You may assume that the `CML` and `SynchVar` structures have been opened, and you are **not** allowed to use the `MakeRPC` structure.

**Question 4.** (14 = 7 × 2 points) For each of the following CML primitives, indicate its type by writing it next to its name:

CML.sendEvt

CML.recvEvt

CML.wrap

CML.guard

CML.withNack

CML.choose

CML.timeOutEvt

**Question 5.** (21 = 7 × (1 + 2) points) For each of the following CML primitives, indicate **both** its type and whether or not it can block when called (circle Y for “can block” or N for “does not block”):

Y N CML.spawn

Y N CML.send

Y N CML.recv

Y N CML.sendPoll

Y N CML.recvPoll

Y N CML.synch

Y N SynchVar.iPut

Y N SynchVar.iGet

Y N SynchVar.mTake

Y N SynchVar.mSwap

Y N Mailbox.send

Y N Mailbox.recv

Y N Multicast.multicast

Y N Multicast.recv

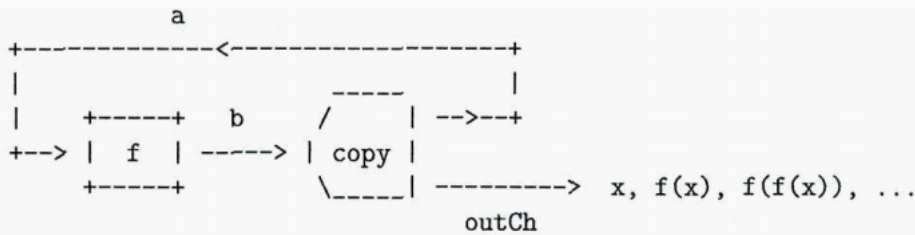
Question 6. (15 points) Recall the process network copy combinator

```
val copy : 'a chan * 'a chan * 'a chan -> unit,
```

where, for channels  $c_1$ ,  $c_2$ , and  $c_3$ , `copy( $c_1, c_2, c_3$ )` creates a process that repeatedly reads values from channel  $c_1$  and writes them to both  $c_2$  and  $c_3$ . For this question, you will use this copy combinator to define a function `iterate` with signature

```
val iterate : 'a -> ('a -> 'a) -> 'a chan
```

The call `iterate  $x$   $f$`  builds the following process network from  $f$  and `copy` and initializes it with the value  $x$ , returning the output channel, `outCh`:



Thus, the return value from the call `iterate  $x$   $f$`  is a *stream* whose values are  $x$ ,  $f(x)$ ,  $f(f(x))$ , and so on. Note that your `iterate` function, besides creating channels for communication and “hooking together” the processes into the network shown and initializing it, will also need to create the *process* labelled “ $f$ ” in the diagram above from the given *function*  $f$ .